

Simplicial Sets within Cubical Sets

Thomas Streicher
TU Darmstadt
jww J. Weinberger

Herrsching December 2019

Computational Meaning of HoTT? (1)

HoTT is **Intensional Type Theory** together with Voevodsky's **Univalence Axiom** (UA) and **Higher Inductive Types** (HITs).

Let us recall what UA says. This requires a few definitions

$$\text{iscontr}(X : \text{Set}) = (\sum x : X)(\prod y : X) \text{Id}_X(x, y)$$

$$\begin{aligned} \text{hfiber}(X, Y : \text{Set})(f : X \rightarrow Y)(y : Y) &= \\ &= (\sum x : X) \text{Id}_Y(f(x), y) \end{aligned}$$

$$\begin{aligned} \text{isweq}(X, Y : \text{Set})(f : X \rightarrow Y) &= \\ &= (\prod y : Y) \text{iscontr}(\text{hfiber}(X, Y, f, y)) \end{aligned}$$

$$\text{Weq}(X, Y : \text{Set}) = (\sum f : X \rightarrow Y) \text{isweq}(X, Y, f)$$

Computational Meaning of HoTT? (2)

Using the eliminator J for identity types one easily defines a map

$$\text{eqweq}(X, Y : \text{Set}) : \text{Id}_{\text{Set}}(X, Y) \rightarrow \text{Weq}(X, Y)$$

Then the **Univalence Axiom**

$$\text{UA} : (\prod X, Y : \text{Set}) \text{isweq}(\text{eqweq}(X, Y))$$

postulates that all maps $\text{eqweq}(X, Y)$ are weak equivalences.

It has been shown that **simplicial sets** provide a model of HoTT interpreting types as Kan complexes.

But UA as it is **lacks computational meaning**:
what should be rewrite rules for the constant UA?

T. Coquand et.al. [CCHM] have developed a *Cubical Type Theory* with computational meaning in which one can **derive** UA.

From Simplicial Sets to Cubical Sets (1)

Bezem and Coquand have shown that the theory of simplicial sets is **not constructive**: one cannot even show constructively that Kan complexes are closed under exponentiation!

This limitation, however, can be overcome when working in **cubical sets**.

The crucial property is that representable objects are closed under finite products: in **sSet** the interval \mathbb{I} is representable but $\mathbb{I} \times \mathbb{I}$ isn't!

The site of **sSet** is Δ , the full subcategory of **Poset** on finite non-empty linear posets.

The site of **cSet** is \square , the full subcategory of **Poset** of finite powers of $\mathbf{2}$.

Splitting idempotents in \square gives rise to **FL**, the full subcategory of **Poset** on finite lattices.

From Simplicial Sets to Cubical Sets (2)

Notice that \square is op-equivalent to the category of *free finitely generated distributive lattices*. Coquand et.al. – just for convenience – used the opposite of the category of *free finitely generated de Morgan algebras*.

Let $i : \Delta \rightarrow \mathbf{FL}$ be the inclusion functor. The restriction functor i^* from $\mathbf{cSet} = \widehat{\mathbf{FL}}$ to $\mathbf{sSet} = \widehat{\Delta}$ has left and right adjoints $i_!$ and i_* , respectively. Since the restriction of the nerve functor Nv to \mathbf{FL} is given by $i^* \circ Y_{\mathbf{FL}}$ we have

$$i_*(X)(L) \cong \mathbf{cSet}(Y_{\mathbf{FL}}(L), i_*(X)) \cong \mathbf{sSet}(i^*Y_{\mathbf{FL}}(L), X) \cong \mathbf{sSet}(Nv(L), X)$$

from which it follows that i_* (and thus also $i_!$) is full and faithful. Since i^* has a left adjoint $i_!$ (given by left Kan extension of $Y_{\mathbf{FL}} \circ i$ along Y_{Δ}) it preserves (finite) limits and thus $i^* \dashv i_*$ is an **injective geometric morphism**.

The topology on **FL** inducing **sSet**

A sieve $S \subseteq Y_{\mathbf{FL}}(L)$ **covers** L iff $i^*S = i^*Y_{\mathbf{FL}}(L) = \text{Nv}(L)$ iff S contains all chains in L , i.e. all monotone maps $[n] \rightarrow L$.

There is a minimal covering sieve C_L consisting of all maps to L whose image is linearly ordered.

The corresponding **closure operator** $j : \Omega \rightarrow \Omega$ sends $S \subseteq Y_{\mathbf{FL}}(L)$ to the set of all $u : K \rightarrow L$ with $uc \in S$ for all chains $c : [n] \rightarrow K$.

Cisinski Model Structures on **cSet** and **sSet** (1)

For Cisinski model structures on presheaf toposes $\widehat{\mathbb{C}}$ its class of **cofibrations** consists of all monos.

Its class of trivial fibrations consists of maps weakly right orthogonal to all monos.

A **naive fibration** is a map weakly right orthogonal to **all cylinders**, i.e. monos of the form

$$(\{\varepsilon\} \times X) \cup (\mathbb{I} \times Y) \hookrightarrow \mathbb{I} \times X$$

where $Y \subseteq X$, $\varepsilon \in \{0, 1\} \subseteq \mathbb{I} = Y(2)$.

The **trivial cofibrations** are the maps weakly left orthogonal to all naive fibrations.

On **sSet** this construction gives the classical model structure and on **cSet** the “type-theoretic” one of [CCHM].

Cisinski Model Structures on **cSet** and **sSet** (2)

Since i^* preserves \mathbb{I} and finite limits and it retracts monos of **cSet** onto monos of **sSet** it follows that i^* retracts cylinders in **sSet** onto cylinders in **sSet**.

Thus i_*p is a fibration in **cSet** iff p is a fibration in **sSet**.

Accordingly, we have $\mathcal{F}_{\mathbf{sSet}} = \mathbf{sSet} \cap \mathcal{F}_{\mathbf{cSet}}$.

Thus $i^* \dashv i_*$ is a *Quillen pair*.

We don't know whether it is a Quillen equivalence.

But presumably not!

Sattler's Theorem

The functor $i_!$ preserves monos, i.e. cofibrations, and also trivial cofibrations.

Thus $i_! \dashv i^*$ is also a Quillen pair.

But we also do not know whether it is a Quillen equivalence, i.e. whether all counits $\varepsilon_X : i_! i^* X \rightarrow X$ are weak equivalences w.r.t. the “type-theoretic” model structure on **cSet**.

i_* preserves and reflects “equality” of fibrant objects

Theorem If A and B are fibrant in **sSet**, i.e. Kan complexes, a map $f : A \rightarrow B$ is a weak equivalence in **sSet** iff it is a weak equivalence in **cSet**.

Proof.

One can show that $i_*(\text{hfiber}(f)) \simeq \text{hfiber}(i_*f)$ and thus

$\forall m \in \text{Mono}(\mathbf{cSet})(m \perp \text{hfiber}(i_*f))$ iff

$\forall m \in \text{Mono}(\mathbf{cSet})(m \perp i_*(\text{hfiber}(f)))$ iff

$\forall m \in \text{Mono}(\mathbf{cSet})(i^*m \perp \text{hfiber}(f))$ iff¹

$\forall m \in \text{Mono}(\mathbf{sSet})(m \perp \text{hfiber}(f))$

i.e.² i_*f is a weak equivalence in **cSet** iff f is a weak equivalence in **sSet**. □

¹since the monos in **sSet** are precisely the sheafifications of monos in **cSet**

²as shown by Voevodsky for fibrant objects A and B a map $w : A \rightarrow B$ is a weak equivalence iff $\text{hfiber}(w)$ is a trivial cofibrations, i.e. $m \perp \text{hfiber}(w)$ for all monos m

Universes in **cSet** and **sSet** (1)

In both **cSet** and **sSet** we may construct universes

$$\pi_c : E_c \rightarrow U_c \quad \text{and} \quad \pi_s : E_s \rightarrow U_s$$

generic for small fibrations.

In both cases the universes are fibrant and univalent!

However, in case of **sSet** this requires heavy choice (due to use of minimal fibrations!) but not so for **cSet**.

Thus, we have

$$\begin{array}{ccccc} E_s & \longrightarrow & i^* E_c & \longrightarrow & E_s \\ \pi_s \downarrow & \lrcorner & i^* \pi_c \downarrow & \lrcorner & \downarrow \pi_s \\ U_s & \xrightarrow{e} & i^* U_c & \xrightarrow{p} & U_s \end{array}$$

and $pe \sim \text{id}_{U_s}$ (i.e. pe and id_{U_s} are homotopy equivalent) since the universe π_s is univalent.

Universes in **cSet** and **sSet** (2)

Since i_* preserves fibrations, pullbacks and \sim for maps between fibrant objects we have

$$\begin{array}{ccccc}
 i_* E_s & \longrightarrow & i_* i^* E_c & \longrightarrow & i_* E_s \\
 i_* \pi_s \downarrow & \lrcorner & i_* i^* \pi_c \downarrow & \lrcorner & \downarrow i_* \pi_s \\
 i_* U_s & \xrightarrow{i_* e} & i_* i^* U_c & \xrightarrow{i_* p} & i_* U_s
 \end{array}$$

with $i_* p \circ i_* e = i_*(p \circ e) \sim i_*(\text{id}_{U_s}) = \text{id}_{i_* U_s}$. Thus $i_* i^* \pi_c$ is generic for small fibrations which are families of sheaves and $i_* \pi_s$ is a univalent such universe.

Thus, pulling back $i_* i^* \pi_c$ along the homotopy equalizer of $i_* e \circ i_* p$ and $\text{id}_{i_* i^* U_c}$ gives rise to a univalent universe U_{cs} in **cSet** which is weakly generic for small fibrations which are families of sheaves.

Conclusion and Outlook

- We have shown that simplicial sets form a(n essential) subtopos of cubical sets.
- Moreover, **sSet** is a submodel of **cSet** since the inclusion $i_* : \mathbf{sSet} \hookrightarrow \mathbf{cSet}$ preserves Σ and Π and also the interval \mathbb{I} and thus also identity types.
- In **cSet** we have constructed a fibrant univalent universe generic for small fibrations of sheaves.

Thus, we conclude that

the simplicial model lives within the cubical model!