

When is the Kleene–Kreisel Hierarchy Full?

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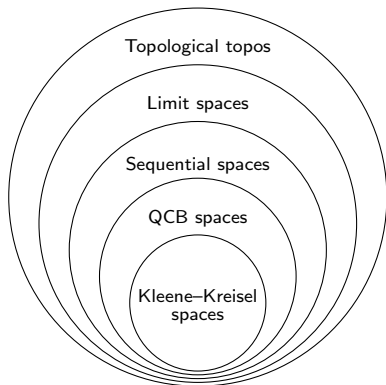
Kleene–Kreisel continuous functionals

First discussed as

- ▶ Kleene's **countable functionals**
- ▶ Kreisel's **continuous functionals**

Calculated within

- ▶ **Compactly generated topological spaces**
- ▶ **Sequential topological spaces**
- ▶ Simpson and Schröders **QCB spaces**
- ▶ Kuratowski **limit spaces**
- ▶ **Filter spaces**
- ▶ Scott's **equilogical spaces**
- ▶ Johnstone's **topological topos**
- ▶ ...



Validating the uniform continuity axiom

The model of Kleene–Kreisel continuous functionals validates the uniform continuity axiom (UC):

$$\forall f: \mathbf{2}^{\mathbb{N}} \rightarrow \mathbb{N}. \exists n \in \mathbb{N}. \forall \alpha, \beta \in \mathbf{2}^{\mathbb{N}}. (\alpha =_n \beta \implies f\alpha = f\beta).$$

But the treatment of the model is non-constructive.

Making it constructive

We introduce a **classically equivalent** model, **C-spaces**,

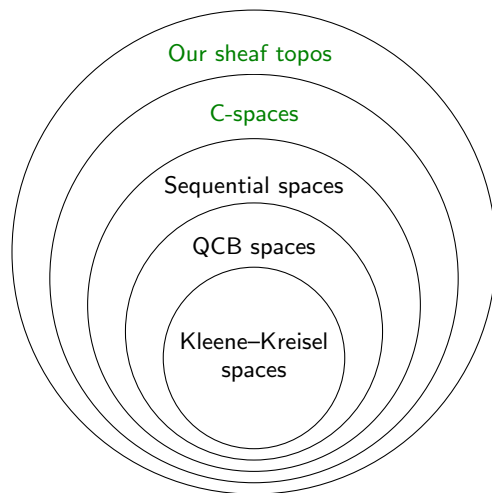
- ▶ in which the uniform continuity axiom holds,
- ▶ but without assuming any constructively contentious axiom in the meta-theory used to define the model.

We work with an intensional type theory with a universe, Σ -types, Π -types, identity types and standard base types.

We have **formalized** our development and proofs in **Agda**.

In this talk, however, I will use **informal**, rigorous mathematical language.

Our constructive manifestation



C-spaces and continuous maps

Def. A **C-topology** on a set X is a collection P of probes $\mathbf{2}^{\mathbb{N}} \rightarrow X$ subject to the following **probe axioms**:

1. All constant maps are in P .
2. If $t: \mathbf{2}^{\mathbb{N}} \rightarrow \mathbf{2}^{\mathbb{N}}$ is uniformly continuous and $p \in P$, then $p \circ t \in P$.
(Presheaf condition)
3. For any two maps $p_0, p_1 \in P$, the unique map $p: \mathbf{2}^{\mathbb{N}} \rightarrow X$ defined by $p(i * \alpha) = p_i(\alpha)$ is in P .
(Sheaf condition)

A **C-space** is a set X equipped with **C-topology**.

A function $f: X \rightarrow Y$ of **C-spaces** is **continuous** if $f \circ p \in P_Y$ whenever $p \in P_X$. (Naturality condition)

We write **C-Space** for the category of **C-spaces** and continuous maps.

Examples of C-spaces

All **continuous** maps from $\mathbf{2}^{\mathbb{N}}$ (with the usual topology) to any topological space X form a C-topology on X :

- ▶ Any constant map $\mathbf{2}^{\mathbb{N}} \rightarrow X$ is continuous.
- ▶ The composite $\mathbf{2}^{\mathbb{N}} \xrightarrow{t} \mathbf{2}^{\mathbb{N}} \xrightarrow{p} X$ of two continuous maps is continuous.
- ▶ The sheaf condition is satisfied because $\mathbf{2}^{\mathbb{N}}$ is compact Hausdorff.

Any continuous map of topological spaces is continuous w.r.t. the above C-topology, as composition preserves continuity.

C-spaces form a (locally) cartesian closed category

The constructions are the same as in the category of sets, with suitable C-topologies. For example,

1. to get products in **C-Space**, we C-topologize cartesian products,
2. to get exponentials in **C-Space**, we C-topologize the sets of continuous maps,
3. to get products in **C-Space**/ X , we C-topologize pullbacks,
4. to get exponentials in **C-Space**/ X , we C-topologize the domains of exponentials in **Set**/ X .

Moreover, **C-Space** has coproducts.

Discrete C-spaces

Def. A map $p: \mathbf{2}^{\mathbb{N}} \rightarrow X$ into a set X is called **locally constant** iff $\exists n \in \mathbb{N}. \forall \alpha, \beta \in \mathbf{2}^{\mathbb{N}}. (\alpha =_n \beta \implies p(\alpha) = p(\beta))$.

Lemma

The locally constant maps $\mathbf{2}^{\mathbb{N}} \rightarrow X$ form a C-topology which has the smallest amount of probes on X .

Def. A C-space X is **discrete** if all functions $X \rightarrow Y$ into any C-space Y are continuous.

Lemma

A C-space is discrete iff its probes are precisely the locally constant functions.

Def. We thus refer to the collection of all locally constant maps $\mathbf{2}^{\mathbb{N}} \rightarrow X$ as the discrete C-topology on X .

Booleans and natural numbers object

The discrete C-topology on $\mathbf{2}$ or \mathbb{N} is the set of uniformly continuous maps.

Theorem

In the category of C-spaces:

1. The discrete space $\mathbf{2}$ is the coproduct of two copies of the terminal space.
2. The discrete space \mathbb{N} is the natural numbers object.

Proof

The unique maps g and h in **Set** in the diagrams below are continuous by the discreteness of $\mathbf{2}$ and \mathbb{N} :

$$\begin{array}{ccccc}
 1 & \xrightarrow{\text{in}_0} & \mathbf{2} & \xleftarrow{\text{in}_1} & 1 \\
 & \searrow g_0 & | g & \swarrow g_1 & \\
 & & X & &
 \end{array}$$

$$\begin{array}{ccccc}
 1 & \xrightarrow{0} & \mathbb{N} & \xrightarrow{\text{suc}} & \mathbb{N} \\
 & \searrow x & | h & & | h \\
 & & X & \xrightarrow{f} & X
 \end{array}$$

Yoneda Lemma

The monoid \mathcal{C} of uniformly continuous $\mathbf{2}^{\mathbb{N}} \rightarrow \mathbf{2}^{\mathbb{N}}$ is a \mathcal{C} -topology on $\mathbf{2}^{\mathbb{N}}$.

The **Yoneda embedding** maps the monoid \mathcal{C} to the \mathcal{C} -space $(\mathbf{2}^{\mathbb{N}}, \mathcal{C})$.

Moreover,

$$y(\star) = (\mathbf{2}^{\mathbb{N}}, \mathcal{C}) = \text{the exponential of the two discrete } \mathcal{C}\text{-spaces.}$$

The **Yoneda Lemma** says that a map $\mathbf{2}^{\mathbb{N}} \rightarrow X$ into a \mathcal{C} -space X is a probe iff it is continuous in the sense of the category of \mathcal{C} -spaces.

The Fan functional

Lemma

The exponential $\mathbb{N}^{2^{\mathbb{N}}}$ is a discrete C-space.

Theorem

There is a continuous functional $\text{fan}: \mathbb{N}^{2^{\mathbb{N}}} \rightarrow \mathbb{N}$ that calculates (minimal) moduli of uniform continuity.

Proof sketch

1. Because any $f \in \mathbb{N}^{2^{\mathbb{N}}}$ is uniformly continuous, we can let $\text{fan}(f)$ be the least witness of this fact.
2. Since $\mathbb{N}^{2^{\mathbb{N}}}$ is discrete according to the above lemma, the functional fan is continuous.

Modelling uniform continuity

C-spaces provide a model of system \mathbf{T} and dependent types:

1. Cartesian closed structure — simply typed λ -calculus.
2. Locally cartesian closed structure — dependent types.
3. Natural numbers object — base type and primitive recursion principle.

Theorem

The uniform continuity axiom is validated by the fan functional.

$$\forall f: \mathbf{2}^{\mathbb{N}} \rightarrow \mathbb{N}. \exists n \in \mathbb{N}. \forall \alpha, \beta \in \mathbf{2}^{\mathbb{N}}. (\alpha =_n \beta \implies f\alpha = f\beta).$$

Kleene–Kreisel functionals via sequence convergence

One way of describing the Kleene–Kreisel spaces is to work with the cartesian closed category of Kuratowski **limit spaces**.

- ▶ A **limit space** is a set together with a designated set of **convergent sequences**, subject to suitable axioms.
- ▶ A function of limit spaces is called **continuous** if it preserves limits.
- ▶ We write **Lim** for the category of limit spaces and continuous maps.
- ▶ To get the Kleene–Kreisel spaces, we start with the discrete natural numbers and iterate exponentials.
- ▶ Example: Any topological space with all (topologically) convergent sequences is a limit space.

C-spaces and limit spaces

- ▶ If A is a limit space, we can give a \mathbf{C} -topology on it by saying that a map $\mathbf{2}^{\mathbb{N}} \rightarrow A$ is a probe on A iff it is limit-continuous.
- ▶ A map is probe-continuous if it is limit-continuous.
- ▶ $G: \mathbf{Lim} \rightarrow \mathbf{C-Space}$
- ▶ If X is a \mathbf{C} -space, we obtain its limit structure by saying that $(x_i) \rightarrow x_\infty$ in X iff the induced function $x: \mathbb{N}_\infty \rightarrow X$ is probe-continuous.
(Uses **non-constructive** arguments.)
- ▶ A map is limit-continuous if it is probe-continuous.
- ▶ $F: \mathbf{C-Space} \rightarrow \mathbf{Lim}$
- ▶ Both F and G keep the underlying set but change the structure, and are identity on morphisms.

Kleene–Kreisel functionals within C-spaces

Lemma

1. $G: \mathbf{Lim} \rightarrow \mathbf{C-Space}$ is a full and faithful embedding.
2. $F: \mathbf{C-Space} \rightarrow \mathbf{Lim}$ is left adjoint to G .

$$\mathbf{Lim} \begin{array}{c} \xleftarrow{F} \\ \perp \\ \xrightarrow{G} \end{array} \mathbf{C-Space}$$

Thus \mathbf{Lim} is a reflective subcategory of $\mathbf{C-Space}$.

3. Moreover, the embedding G preserves the natural numbers object, and the cartesian closed structure (products and exponentials).

Theorem

Kleene–Kreisel continuous functionals can be calculated within $\mathbf{C-Space}$.

If UC holds in the meta-theory

What does our model construction do?

Def.

The collection of all maps $\mathbf{2}^{\mathbb{N}} \rightarrow X$ form an **indiscrete C-topology** on X .

We say X is an **indiscrete C-space**.

Lemma

1. The category of indiscrete C-spaces is equivalent to **Set**.
2. Indiscrete C-spaces form an exponential ideal of **C-Space**.
3. If **UC** holds in **Set**, then the discrete space \mathbb{N} is also indiscrete.

The Kleene–Kreisel and full type hierarchies

Def.

Let \mathbf{C} be a cartesian closed category with a natural numbers object. The **type hierarchy** on \mathbf{C} is the smallest full subcategory containing the natural numbers object and closed under exponentials.

The one on \mathbf{Set} is called the **full type hierarchy**.

The one on $\mathbf{C-Space}$ is called the **Kleene–Kreisel hierarchy**.

Corollary

If **UC** holds in \mathbf{Set} , then

the Kleene–Kreisel hierarchy is equivalent to the full type hierarchy.

Theorem

The Kleene–Kreisel hierarchy and the full type hierarchy are equivalent **if and only if UC** holds in \mathbf{Set} .

Summary

1. Validation of uniform continuity axiom in a weak constructive meta-theory.
2. Constructive manifestation of Kleene–Kreisel continuous functionals.
3. Equivalence of the Kleene–Kreisel hierarchy and the full type hierarchy when assuming **UC**.
4. Compatible with intuitionistic type theory and formalized in Agda.
5. Extraction of computational content from type-theoretic proofs which use **UC**.